

Fully Invariant and Verbal Congruences

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Verbal Congruences

A an algebra, Σ a set of equations

$$\lambda_{\Sigma}^{\mathbf{A}} = \text{Cg}^{\mathbf{A}} \left\{ (s(a_1, \dots, a_n), t(a_1, \dots, a_n)) : \right. \\ \left. (s \approx t) \in \Sigma, a_1, \dots, a_n \in \mathbf{A} \right\}$$

$\lambda_{\Sigma}^{\mathbf{A}}$ is the smallest congruence θ such that $\mathbf{A}/\theta \models \Sigma$

$\lambda_{\Sigma}^{\mathbf{A}}$ is the *verbal congruence induced by Σ*

Alternate definition

\mathcal{V} a variety (same similarity type as \mathbf{A})

$$\Lambda_{\mathcal{V}}^{\mathbf{A}} = \{ \theta \in \text{Con}(\mathbf{A}) : \mathbf{A}/\theta \in \mathcal{V} \}$$

$$\lambda_{\mathcal{V}}^{\mathbf{A}} = \bigcap \Lambda_{\mathcal{V}}^{\mathbf{A}}$$

Easy to see:

- $\mathbf{A}/\lambda_{\mathcal{V}} \in \mathcal{V}$
- If $\mathcal{V} = \text{Mod}(\Sigma)$ then $\lambda_{\mathcal{V}} = \lambda_{\Sigma}$

Theorem

θ is verbal on \mathbf{A} iff

$$\mathbf{A}/\psi \in \text{Var}(\mathbf{A}/\theta) \implies \psi \geq \theta.$$

Examples from Group Theory

Suppose $\Sigma = \{xy \approx yx\}$

On any group \mathbf{A} , $\lambda_{\Sigma}^{\mathbf{A}} = \text{Cg} \{ (ab, ba) : a, b \in A \}$
corresponds to $\mathbf{A}' = \text{Nml} \{ [a, b] : a, b \in A \} = [A, A]$

\mathbf{A}/\mathbf{A}' is the largest Abelian homomorphic image of \mathbf{A} .

$$\Theta_n = \{x^n \approx e\}$$

$\lambda_{\Theta_n}^{\mathbf{A}}$ corresponds to $\text{Nml} \{ a^n : a \in A \}$

$\mathbf{A}/\lambda_{\Theta_n}$ is the largest homomorphic image of \mathbf{A} of exponent n .

Note: $\lambda_{\Sigma} \leq \lambda_{\Theta_2}$ since every group of exponent 2 is Abelian

Fully Invariant Congruences

$\text{End}(\mathbf{A})$ = endomorphism monoid of \mathbf{A}

A congruence θ is *fully invariant* if

$$\forall f \in \text{End}(\mathbf{A}) \quad (a, b) \in \theta \implies (f(a), f(b)) \in \theta$$

Theorem

Every verbal congruence is fully invariant

Converse is false

Example: Let p be prime,
 \mathbf{A} the Abelian group $\langle a, b \mid pa = p^2b = 0 \rangle \cong \mathbb{Z}_p \oplus \mathbb{Z}_{p^2}$.

$$S = \{x \in \mathbf{A} : px = 0\} \supsetneq T = \{px : x \in \mathbf{A}\}$$

S is fully invariant but

$\mathbf{A}/T \in \text{Var}(\mathbf{A}/S)$ so S not verbal

General Question

Find conditions under which

$$\text{fully invariant} \implies \text{verbal}$$

- for a congruence
- for an algebra
- for a variety.

An algebra is called *verbose* if every fully invariant congruence is verbal.

A variety is *verbose* if every member is verbose.

Verbose Varieties

Theorem

- *A variety of Abelian groups is verbose if and only if it is of square-free exponent.*
- *The lattice varieties $\mathbf{V}(\mathbf{M}_n, \mathbf{N}_5)$ are verbose. The variety $\mathbf{V}(\mathbf{M}_{3,3})$ is not verbose.*

Theorem

Let \mathcal{V} be a finitely generated discriminator variety. Then \mathcal{V}_{fin} is verbose.

Proof.

$\mathbf{A} \in \mathcal{V}_{fin} \implies \mathbf{A} \cong \mathbf{A}_1 \times \cdots \times \mathbf{A}_n$ all simple.

$\eta_i = \ker(\mathbf{A} \twoheadrightarrow \mathbf{A}_i)$.

Suppose $\theta \in \text{Con}(\mathbf{A})$ not verbal. Then

$\exists i, j \ \theta \leq \eta_i, \theta \not\leq \eta_j, \mathbf{A}_j \xrightarrow{h} \mathbf{A}_i$

Define $e(\mathbf{x}) = (x_1, \dots, x_{i-1}, h(x_j), x_{i+1}, \dots, x_n)$.

Then $e \in \text{End}(\mathbf{A})$ but

$(\mathbf{a}, \mathbf{b}) \in \theta - \eta_j \implies (e(\mathbf{a}), e(\mathbf{b})) \notin \theta$.



Can this argument be extended to infinite algebras?

Need a representation that is “almost as good” as direct product

Answer: NU-duality

A finite algebra \mathbf{M} is *subalgebra-primal* if $\text{Clo}(\mathbf{M}) = \text{Pol}(\text{Sub}(\mathbf{M}))$.

\mathbf{M} is quasi-primal and rigid

Assume $\mathcal{V} = \text{Var}(\mathbf{M})$, \mathbf{M} subalgebra-primal

$\mathbf{A} \in \mathcal{V} \implies \mathbf{A} \rightsquigarrow \langle X, T, S \rangle$, T a Boolean Topology on X , S a family of closed subspaces

Instead of constructing endomorphism $\mathbf{A} \xrightarrow{e} \mathbf{A}$
 build continuous map $X \xleftarrow{\hat{e}} X$ preserving S

We can imitate the finite argument to show \mathcal{V} is verbose

Should be true for finitely generated discriminator variety.
Need the continuous functions $X \xleftarrow{\widehat{\sigma}} X$ to preserve the
“inner automorphisms” of \mathbf{A} .

Middle ground: Can show that $\mathbf{V}(\mathbf{F})$ is verbal where \mathbf{F} is a
finite field.