

Topological systems versus attachment relations

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INVESTMENTS IN EDUCATION DEVELOPMENT

Outline

- 1 Introduction: topological systems and attachment relations
- 2 Variety-based topological systems
- 3 Variety-based attachment relations
- 4 Topological systems versus attachment relations
- 5 Conclusion: open problem

Topological systems and attachment relations

- There exist two notions related to topology: *topological system* of S. Vickers and *attachment relation* of C. Guido.
- Topological systems were introduced in 1989 as a common framework for both topological spaces and their underlying algebraic structures – frames or locales.
- Attachment relations were introduced in 2009 as a fuzzification of the relation “ \in ” (“element of”) between points and sets.

Variety-based approach

- In 2010, the notions of *variety-based topological system* and *attachment relation* were introduced, thereby suggesting the problem of comparing systems and attachments.
- Both concepts are given by an indexed family of homomorphisms of algebras of some variety, where topological systems use an indexing over a set, and attachment relations put the algebraic structure of the variety in question on it.

Topological systems versus attachment relations

- This talk shows an embedding of the category of attachment relations into the category of topological systems.
- We also construct a functor in the opposite direction, which is an embedding under certain restrictions on its domain.
- The two functors make a link between the spatialization and localization procedures for systems and attachments, where the attachment localization procedure is defined for this occasion.
- We construct, additionally, another functor from systems to attachments, and provide a natural transformation between the two system-to-attachment ways.

Ω -algebras and Ω -homomorphisms

Definition 1

Let $\Omega = (n_\lambda)_{\lambda \in \Lambda}$ be a class of cardinal numbers.

- An **Ω -algebra** is a pair $(A, (\omega_\lambda^A)_{\lambda \in \Lambda})$, comprising a set A and a family of maps $A^{n_\lambda} \xrightarrow{\omega_\lambda^A} A$ (**n_λ -ary primitive operations** on A).
- An **Ω -homomorphism** $(A_1, (\omega_\lambda^{A_1})_{\lambda \in \Lambda}) \xrightarrow{\varphi} (A_2, (\omega_\lambda^{A_2})_{\lambda \in \Lambda})$ is a map $A_1 \xrightarrow{\varphi} A_2$ such that $\varphi \circ \omega_\lambda^{A_1} = \omega_\lambda^{A_2} \circ \varphi^{n_\lambda}$ for every $\lambda \in \Lambda$.
- **$\text{Alg}(\Omega)$** is the construct of Ω -algebras and Ω -homomorphisms.

Varieties and algebras

Definition 2

Let \mathcal{M} (resp. \mathcal{E}) be the class of Ω -homomorphisms with injective (resp. surjective) underlying maps. A *variety of Ω -algebras* is a full subcategory of $\mathbf{Alg}(\Omega)$, which is closed under the formation of products, \mathcal{M} -subobjects and \mathcal{E} -quotients, and whose objects (resp. morphisms) are called *algebras* (resp. *homomorphisms*).

Definition 3

Given a variety \mathbf{A} , a *reduct* of \mathbf{A} is a pair $(|-|, \mathbf{B})$, where \mathbf{B} is a variety such that $\Omega_{\mathbf{B}} \subseteq \Omega_{\mathbf{A}}$ (every primitive operation of \mathbf{B} is a primitive operation of \mathbf{A}), whereas $\mathbf{A} \xrightarrow{|-|} \mathbf{B}$ is a concrete functor.

Variety-based powerset operators

Theorem 4

Given a variety \mathbf{A} , every subcategory \mathbf{S} of \mathbf{A}^{op} induces the functor $\mathbf{Set} \times \mathbf{S} \xrightarrow{(-)^{\leftarrow}} \mathbf{A}^{op}$, which is defined by $(X_1, A_1) \xrightarrow{(f, \varphi)^{\leftarrow}} (X_2, A_2) = A_1^{X_1} \xrightarrow{((f, \varphi)^{\leftarrow})^{op}} A_2^{X_2}$, where $(f, \varphi)^{\leftarrow}(\alpha) = \varphi^{op} \circ \alpha \circ f$.

Example 5

- ① If $\mathbf{A} = \mathbf{CBAAlg}$ (complete Boolean algebras) and $\mathbf{S} = \{2 \xrightarrow{1_2} 2\}$, then $\mathbf{Set} \times \mathbf{S} \xrightarrow{(-)^{\leftarrow}} \mathbf{CBAAlg}^{op}$ is the classical preimage operator.
- ② If $\mathbf{A} = \mathbf{Frm}$ (frames), then $\mathbf{Set} \times \mathbf{S} \xrightarrow{(-)^{\leftarrow}} (\mathbf{Loc} := \mathbf{Frm}^{op})$ (locales) is the powerset operator of \mathbf{S} . E. Rodabaugh.

Variety-based topological spaces

Definition 6

Given a variety \mathbf{A} , a subcategory \mathbf{S} of \mathbf{A}^{op} and a reduct \mathbf{B} of \mathbf{A} , $\mathbf{S}_{\mathbf{B}}\text{-Top}$ is the category, concrete over the category $\mathbf{Set} \times \mathbf{S}$, which comprises the following data.

Objects: triples (X, A, τ) ($\mathbf{S}_{\mathbf{B}}\text{-topological spaces}$ or $\mathbf{S}_{\mathbf{B}}\text{-spaces}$), where (X, A) is a $\mathbf{Set} \times \mathbf{S}$ -object, and τ ($\mathbf{S}_{\mathbf{B}}\text{-topology}$ on (X, A)) is a \mathbf{B} -subalgebra of $|A^X|$.

Morphisms: $(X_1, A_1, \tau_1) \xrightarrow{(f, \varphi)} (X_2, A_2, \tau_2)$ are $\mathbf{Set} \times \mathbf{S}$ -morphisms $(X_1, A_1) \xrightarrow{(f, \varphi)} (X_2, A_2)$ with the property that $(f, \varphi)^{\leftarrow}(\alpha_2) \in \tau_1$ for every $\alpha_2 \in \tau_2$ ($\mathbf{S}_{\mathbf{B}}\text{-continuity}$).

Examples

Example 7

- ① If $\mathbf{A} = \mathbf{CBAAlg}$, $\mathbf{S} = \{2 \xrightarrow{1_2} 2\}$, and $\mathbf{B} = \mathbf{Frm}$, then $\mathbf{S}_\mathbf{B}\text{-Top}$ is the category \mathbf{Top} of topological spaces.
- ② If $\mathbf{A} = \mathbf{B} = \mathbf{Frm}$, then $\mathbf{S}_\mathbf{B}\text{-Top}$ is the category $\mathbf{S}\text{-Top}$ of lattice-valued topological spaces of \mathbf{S} . E. Rodabaugh.

Variety-based topological systems

Definition 8

Given a variety \mathbf{A} , a subcategory \mathbf{S} of \mathbf{A}^{op} and a reduct \mathbf{B} of \mathbf{A} , $\mathbf{S}_{\mathbf{B}}\text{-TopSys}$ is the category, concrete over the category $\mathbf{Set} \times \mathbf{B}^{op} \times \mathbf{S}$, which comprises the following data.

Objects: tuples $D = (\text{pt } D, \Omega D, \Sigma D, \models)$ ($\mathbf{S}_{\mathbf{B}}\text{-topological systems}$ or $\mathbf{S}_{\mathbf{B}}\text{-systems}$), where $(\text{pt } D, \Omega D, \Sigma D)$ is a $\mathbf{Set} \times \mathbf{B}^{op} \times \mathbf{S}$ -object, and $\text{pt } D \times \Omega D \xrightarrow{\models} \Sigma D$ is a map ($\Sigma D\text{-satisfaction relation}$) such that $\Omega D \xrightarrow{\models(x, -)} |\Sigma D|$ is a \mathbf{B} -homomorphism for every $x \in \text{pt } D$.

Morphisms: $D_1 \xrightarrow{f = (\text{pt } f, (\Omega f)^{op}, (\Sigma f)^{op})} D_2$ are $\mathbf{Set} \times \mathbf{B}^{op} \times \mathbf{S}$ -morphisms $(\text{pt } D_1, \Omega D_1, \Sigma D_1) \xrightarrow{f} (\text{pt } D_2, \Omega D_2, \Sigma D_2)$ such that $\Sigma f \circ \models_2(\text{pt } f(x), b) = \models_1(x, \Omega f(b))$ for every $x \in \text{pt } D_1, b \in \Omega D_2$.

Examples

Example 9

- ① If $\mathbf{A} = \mathbf{CBAAlg}$, $\mathbf{B} = \mathbf{Frm}$, and $\mathbf{S} = \{2 \xrightarrow{1_2} 2\}$, then $\mathbf{S}_B\text{-TopSys}$ is the category **TopSys** of topological systems of S . Vickers.
- ② If $\mathbf{A} = \mathbf{B} = \mathbf{Frm}$, and $\mathbf{S} = \mathbf{Loc}$, then $\mathbf{S}_B\text{-TopSys}$ is the category of lattice-valued topological systems of S . E. Rodabaugh.

System spatialization procedure

Theorem 10

- ① *There exists the full embedding $\mathbf{S}_B\text{-Top} \xrightarrow{E_{TS}} \mathbf{S}_B\text{-TopSys}$, which is defined by $E_{TS}((X_1, A_1, \tau_1) \xrightarrow{(f, \varphi)} (X_2, A_2, \tau_2)) = (X_1, \tau_1, A_1, \models_1) \xrightarrow{(f, ((f, \varphi)^\leftarrow |_{\tau_2}^{\tau_1})^{op}, \varphi)} (X_2, \tau_2, A_2, \models_2)$, where $\models_i(x, \alpha) = \alpha(x)$.*
- ② *There exists the functor $\mathbf{S}_B\text{-TopSys} \xrightarrow{\text{Spat}_{TS}} \mathbf{S}_B\text{-Top}$, which is defined by $\text{Spat}_{TS}(D_1 \xrightarrow{f} D_2) = (\text{pt } D_1, \Sigma D_1, \tau_1) \xrightarrow{(\text{pt } f, (\Sigma f)^{op})} (\text{pt } D_2, \Sigma D_2, \tau_2)$, where $\tau_i = \{\models_i(-, b) \mid b \in \Omega D_i\}$.*
- ③ *Spat_{TS} is a right-adjoint-left-inverse to E_{TS} .*

System localification procedure

Theorem 11

- Every object A of the category \mathbf{S} gives rise to an embedding $\mathbf{B}^{op} \xrightarrow{E_{TS}^A} \mathbf{S}_B\text{-TopSys}$ defined by $E_{TS}^A(B_1 \xrightarrow{\varphi} B_2) = (Pt_A(B_1), B_1, A, \models_1) \xrightarrow{((\varphi^{op})_A^{\leftarrow}, \varphi, 1_A)} (Pt_A(B_2), B_2, A, \models_2)$, where $Pt_A(B_i) = \mathbf{B}(B_i, |A|)$ and $\models_i(p, b) = p(b)$. E_{TS}^A is full iff $\mathbf{S}(A, A) = \{1_A\}$.
- There exists the functor $\mathbf{S}_B\text{-TopSys} \xrightarrow{\text{Loc}_{TS}^\Omega} \mathbf{B}^{op}$ defined by $\text{Loc}_{TS}^\Omega(D_1 \xrightarrow{f} D_2) = \Omega D_1 \xrightarrow{(\Omega f)^{op}} \Omega D_2$.
- Loc_{TS}^Ω is a left inverse to E_{TS}^A . If $\mathbf{S} = \{A \xrightarrow{1_A} A\}$, then Loc_{TS}^Ω is a left adjoint to E_{TS}^A .

Variety-based attachment relations

Definition 12

Given a variety \mathbf{A} , a subcategory \mathbf{S} of \mathbf{A}^{op} , a reduct \mathbf{B} of \mathbf{A} , and a functor $\mathbf{B}^{op} \xrightarrow{(-)^*} \mathbf{Set}$ such that $B^* = |B|$, $\mathbf{S}_B\text{-Att}$ is the category, concrete over the category $\mathbf{B}^{op} \times \mathbf{S}$, comprising the following data.

Objects: triples $F = (\Omega F, \Sigma F, \Vdash)$ (**\mathbf{S}_B -attachments**), where $(\Omega F, \Sigma F)$ is a $\mathbf{B}^{op} \times \mathbf{S}$ -object, and $\Omega F \xrightarrow{\Vdash} \mathbf{B}(\Omega F, |\Sigma F|)$ is a map (**ΣF -attachment relation**).

Morphisms: $F_1 \xrightarrow{f = ((\Omega f)^{op}, (\Sigma f)^{op})} F_2$ are $\mathbf{B}^{op} \times \mathbf{S}$ -morphisms $(\Omega F_1, \Sigma F_1) \xrightarrow{f} (\Omega F_2, \Sigma F_2)$ such that $\Sigma f \circ \Vdash_2((\Omega f)^{op*}(b_1))(b_2) = (\Vdash_1(b_1))(\Omega f(b_2))$ for every $b_1 \in \Omega F_1, b_2 \in \Omega F_2$ (**\mathbf{S}_B -continuity**).

Example

Example 13

If $\mathbf{A} = \mathbf{CBA}lg$, $\mathbf{B} = \mathbf{Frm}$, and $\mathbf{S} = \{2 \xrightarrow{1_2} 2\}$, then $\mathbf{S-Att}$ is the category of attachment relations of \mathbf{C} . Guido.

Attachment spatialization procedure

Theorem 14

- 1 There exists the functor $\mathbf{S}_B\text{-Top} \xrightarrow{E_{AT}} \mathbf{S}_B\text{-Att}$, which is given by $E_{AT}((X_1, A_1, \tau_1) \xrightarrow{(f, \varphi)} (X_2, A_2, \tau_2)) = (\tau_1, A_1^{X_1}, \llbracket \tau_1 \rrbracket) \xrightarrow{(((f, \varphi)^{\leftarrow} \llbracket \tau_1 \rrbracket)^{op}, ((f, \varphi)^{\leftarrow})^{op})} (\tau_2, A_2^{X_2}, \llbracket \tau_2 \rrbracket)$, where $(\llbracket \tau_i \rrbracket(\alpha))(\alpha') = \alpha'$. E_{AT} is injective on objects, but, in general, is not faithful and, therefore, fails to be an embedding.
- 2 There exists the functor $\mathbf{S}_B\text{-Att} \xrightarrow{\text{Spat}_{AT}} \mathbf{S}_B\text{-Top}$, which is defined by the formula $\text{Spat}_{AT}(F_1 \xrightarrow{f} F_2) = (|\Omega F_1|, \Sigma F_1, \tau_1) \xrightarrow{((\Omega f)^{op*}, (\Sigma f)^{op})} (|\Omega F_2|, \Sigma F_2, \tau_2)$, where $\tau_i = \{(\llbracket \tau_i \rrbracket(-))(b) \mid b \in \Omega F_i\}$.
- 3 Spat_{AT} is not a left inverse to E_{AT} .

Remarks

While comparing the spatialization procedures for systems and attachments, one could emphasize three important differences.

- The functor E_{AT} is not an embedding.
- Spat_{AT} is not a left inverse to E_{AT} .
- Since the definition of the category **S_B-Att** involves a functor $\mathbf{B}^{op} \xrightarrow{(-)^*} \mathbf{Set}$, whose nature, in general, is unclear, we are unable to determine whether Spat_{AT} is a right adjoint to E_{AT} .

Attachment localization procedure

Theorem 15

- 1 There exists the full embedding $\mathbf{S} \xrightarrow{E_{AT}^S} \mathbf{S}_B\text{-Att}$, which is defined by $E_{AT}^S(A_1 \xrightarrow{\varphi} A_2) = (|A_1|, A_1, \Vdash_1) \xrightarrow{(\varphi, \varphi)} (|A_2|, A_2, \Vdash_2)$, where $\Vdash_i(a) = 1_{A_i}$.
- 2 There exists the functor $\mathbf{S}_B\text{-Att} \xrightarrow{\text{Loc}_{AT}^\Sigma} \mathbf{S}$, which is given by $\text{Loc}_{AT}^\Sigma(F_1 \xrightarrow{f} F_2) = \Sigma F_1 \xrightarrow{(\Sigma f)^{op}} \Sigma F_2$.
- 3 Loc_{AT}^Σ is a left inverse to E_{AT}^S . In general, Loc_{AT}^Σ is not a left adjoint to E_{AT}^S .

Embedding attachments into systems

Theorem 16

There exists the embedding $\mathbf{S}_B\text{-Att} \xrightarrow{H} \mathbf{S}_B\text{-TopSys}$ given by

$$H(F_1 \xrightarrow{f} F_2) =$$

$$(|\Omega F_1|, \Omega F_1, \Sigma F_1, \models_1) \xrightarrow{((\Omega f)^{op*}, (\Omega f)^{op}, (\Sigma f)^{op})} (|\Omega F_2|, \Omega F_2, \Sigma F_2, \models_2),$$

where $\models_i(b, b') = (\Vdash_i(b))(b')$.

From systems to attachments

Theorem 17

There exists the functor $\mathbf{S}_B\text{-TopSys} \xrightarrow{K} \mathbf{S}_B\text{-Att}$ given by

$$K(D_1 \xrightarrow{f} D_2) = (\Omega D_1, (\Sigma D_1)^{\text{pt } D_1}, \Vdash_1) \xrightarrow{((\Omega f)^{\text{op}}, (\text{pt } f, (\Sigma f)^{\text{op}})^{\leftarrow})^{\text{op}}} (\Omega D_2, (\Sigma D_2)^{\text{pt } D_2}, \Vdash_2),$$

where $((\Vdash_i(b))(b'))(x) = \Vdash_i(x, b')$ for every $b, b' \in \Omega D_i, x \in \text{pt } D_i$.
The functor K is injective on objects.

Embedding systems into attachments

Definition 18

$\mathbf{S_B-TopSys}_\Sigma^*$ is the non-full subcategory of $\mathbf{S_B-TopSys}$, comprising the following data.

Objects: $\mathbf{S_B}$ -systems D such that $\text{pt } D$ is non-empty, and ΣD has at least two elements.

Morphisms: $\mathbf{S_B}$ -system morphisms $D_1 \xrightarrow{f} D_2$ such that the map $\Sigma D_2 \xrightarrow{\Sigma f} \Sigma D_1$ is injective.

Theorem 19

The restriction of the functor $\mathbf{S_B-TopSys} \xrightarrow{K} \mathbf{S_B-Att}$ to $\mathbf{S_B-TopSys}_\Sigma^*$ is an embedding.

From systems to attachments again

Theorem 20

- ① *There exists the functor $\mathbf{S}_B\text{-TopSys} \xrightarrow{L} \mathbf{S}_B\text{-Att}$ given by*

$$L(D_1 \xrightarrow{f} D_2) = ((\Omega D_1)^{\text{pt } D_1}, (\Sigma D_1)^{\text{pt } D_1}, \Vdash_1) \xrightarrow{(((\text{pt } f, (\Omega f)^{\text{op}})^{\leftarrow})^{\text{op}}, ((\text{pt } f, (\Sigma f)^{\text{op}})^{\leftarrow})^{\text{op}})} ((\Omega D_2)^{\text{pt } D_2}, (\Sigma D_2)^{\text{pt } D_2}, \Vdash_2),$$

where $((\Vdash_i(\alpha))(\alpha'))(x) = \Vdash_i(x, \alpha'(x))$ for every $\alpha, \alpha' \in (\Omega D_i)^{\text{pt } D_i}$, $x \in \text{pt } D_i$. The functor L is injective on objects.

- ② *The restriction of L to $\mathbf{S}_B\text{-TopSys}_\Sigma^*$ is an embedding.*

Embedding systems into attachments again

Definition 21

$\mathbf{S_B-TopSys}_\Omega^*$ is the non-full subcategory of $\mathbf{S_B-TopSys}$, comprising the following data.

Objects: $\mathbf{S_B}$ -systems D such that $\text{pt } D$ is non-empty, and ΩD has at least two elements.

Morphisms: $\mathbf{S_B}$ -system morphisms $D_1 \xrightarrow{f} D_2$ such that the map $\Omega D_2 \xrightarrow{\Omega f} \Omega D_1$ is injective.

Theorem 22

The restriction of the functor $\mathbf{S_B-TopSys} \xrightarrow{L} \mathbf{S_B-Att}$ to $\mathbf{S_B-TopSys}_\Omega^*$ is an embedding.

Relationships between two embeddings

Theorem 23

There exists a natural transformation

$$\begin{array}{ccc}
 & L & \\
 \text{S}_{\mathbf{B}}\text{-TopSys} & \xrightarrow{\quad} & \text{S}_{\mathbf{B}}\text{-Att,} \\
 & \Downarrow \eta & \\
 & K &
 \end{array}$$

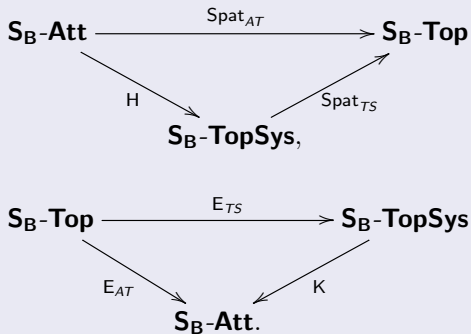
which is defined by the formula $L D \xrightarrow{\eta_D} K D = ((\Omega D)^{\text{pt } D}, (\Sigma D)^{\text{pt } D}, \Vdash_L) \xrightarrow{(([\mathbf{1}_{\Omega D}]_{\text{pt } D})^{\text{op}}, \mathbf{1}_{(\Sigma D)^{\text{pt } D}})} (\Omega D, (\Sigma D)^{\text{pt } D}, \Vdash_K)$.

η provides a connection between *variety-based point-set* (involving powersets $(\Omega D)^{\text{pt } D}$ or $(\Sigma D)^{\text{pt } D}$) and *variety-based pointless* (involving just a \mathbf{B} -algebra ΩD or an \mathbf{A} -algebra ΣD) topologies.

The functors H , K and the spatialization procedure

Theorem 24

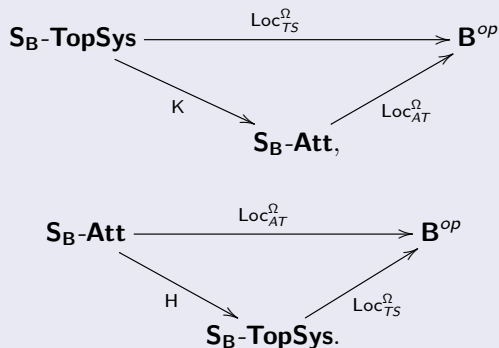
The following triangles commute:



The functors H , K and the localification procedure

Theorem 25

The following triangles commute:



Another system localification procedure

Theorem 26

- ① *There exists the embedding*

$$\mathbf{S} \xrightarrow{E_{TS}^S} \mathbf{S}_B\text{-TopSys} = \mathbf{S} \xrightarrow{E_{AT}^S} \mathbf{S}_B\text{-Att} \xrightarrow{H} \mathbf{S}_B\text{-TopSys}$$

defined by $E_{TS}^S(A_1 \xrightarrow{\varphi} A_2) = (|A_1|, |A_1|, A_1, \models_1) \xrightarrow{(\varphi^*, \varphi, \varphi)} (|A_2|, |A_2|, A_2, \models_2)$, where $\models_i(a, -) = 1_{A_i}$.

- ② *There exists the functor $\mathbf{S}_B\text{-TopSys} \xrightarrow{\text{Loc}_{TS}^\Sigma} \mathbf{S}$, which is given by $\text{Loc}_{TS}^\Sigma(D_1 \xrightarrow{f} D_2) = \Sigma D_1 \xrightarrow{(\Sigma f)^{op}} \Sigma D_2$.*
- ③ *E_{TS}^S is a right inverse to Loc_{TS}^Σ , and, in general, is not full.*

Final remarks

- This talk considered functorial relationships between the categories of (variety-based) topological systems of S. Vickers and attachment relations of C.Guido.
- The talk showed an embedding of the category of attachment relations into the category of topological systems, suggesting a meta-mathematical conclusion that the idea of topological system is more general than the idea of attachment relation.
- We also constructed two functors in the opposite direction, which are embeddings under some restrictions on their domains only, and showed a link between them, which (ultimately) provided two morphisms of variety-based topological theories.





Open problem

This talk constructed an embedding $\mathbf{S}_B\text{-Att} \xrightarrow{H} \mathbf{S}_B\text{-TopSys}$,
 and, additionally, two functors $\mathbf{S}_B\text{-TopSys} \begin{matrix} \xrightarrow{K} \\ \xrightarrow{L} \end{matrix} \mathbf{S}_B\text{-Att}$.

Problem 27

Is it true that the category $\mathbf{S}_B\text{-Att}$ is isomorphic to a (co)reflective subcategory of the category $\mathbf{S}_B\text{-TopSys}$? More precisely, is it true that the embedding H has a left or a right adjoint?

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Thank you for your attention!